

CS103
FALL 2025



Lecture 05:

First-Order Logic

Part 2 of 2

Recap from Last Time

What is First-Order Logic?

- ***First-order logic*** is a logical system for reasoning about properties of objects.
- Augments the logical connectives from propositional logic with
 - ***predicates*** that describe properties of objects,
 - ***functions*** that map objects to one another, and
 - ***quantifiers*** that allow us to reason about many objects at once.

Some bear is curious.

$\exists b. (Bear(b) \wedge Curious(b))$

$\exists$ is the **existential quantifier**
and says "there is a choice of
b where the following is true."

“For any natural number n ,
 n is even if and only if n^2 is even”

$\forall n. (n \in \mathbb{N} \rightarrow (Even(n) \leftrightarrow Even(n^2)))$



$\forall$ is the **universal quantifier**
and says “for any choice of n ,
the following is true.”

“Some P is a Q ”

translates as

$\exists x. (P(x) \wedge Q(x))$

Useful Intuition:

Existentially-quantified statements are false unless there's a positive example.

$$\exists x. (P(x) \wedge Q(x))$$

If x is an example, it must have property P on top of property Q .

“All P 's are Q 's”

translates as

$$\forall x. (P(x) \rightarrow Q(x))$$

Useful Intuition:

Universally-quantified statements are true unless there's a counterexample.

$$\forall x. (P(x) \rightarrow Q(x))$$

If x is a counterexample, it must have property P but not have property Q .

New Stuff!

The Aristotelian Forms

“All As are Bs”

$\forall x. (A(x) \rightarrow B(x))$

“Some As are Bs”

$\exists x. (A(x) \wedge B(x))$

“No As are Bs”

$\forall x. (A(x) \rightarrow \neg B(x))$

“Some As aren’t Bs”

$\exists x. (A(x) \wedge \neg B(x))$

It is worth committing these patterns to memory. We'll be using them throughout the day and they form the backbone of many first-order logic translations.

The Art of Translation

Using the predicates

- *Person*(p), which states that p is a person, and
- *Loves*(x, y), which states that x loves y ,

write a sentence in first-order logic that means “every person loves someone else.”

Answer at

<https://cs103.stanford.edu/pollev>

Every person loves someone else

Every person loves some other person

Every person p loves some other person

Every person p loves some other person

“All A s are B s”

$\forall x. (A(x) \rightarrow B(x))$

$\forall p. (Person(p) \rightarrow$
 p loves some other person

)

“All As are Bs”

$\forall x. (A(x) \rightarrow B(x))$

$\forall p. (Person(p) \rightarrow$
 p loves some other person

)

$\forall p. (Person(p) \rightarrow$
there is some other person that p loves

)

$\forall p. (Person(p) \rightarrow$
there is a person other than p that p loves

)

$\forall p. (Person(p) \rightarrow$
there is a person q , other than p , where p loves q
)

$\forall p. (Person(p) \rightarrow$
 there is a person q , other than p , where
 p loves q
 $)$

$\forall p. (Person(p) \rightarrow$
there is a person q , other than p , where
 p loves q

)

“Some A s are B s”

$\exists x. (A(x) \wedge B(x))$

$\forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge , \text{other than } p, \text{ where}$
 $\text{ } p \text{ loves } q$
 $)$
 $)$

“Some As are Bs”

$\exists x. (A(x) \wedge B(x))$

$\forall p. (Person(p) \rightarrow$
 $\exists q. (Person(q) \wedge$, *other than p, where*
 p loves q
)
)

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. (Person(q) \wedge p \neq q \wedge$$
$$\quad \quad p \text{ loves } q$$
$$\quad)$$
$$)$$

$$\forall p. (Person(p) \rightarrow$$
$$\quad \exists q. (Person(q) \wedge p \neq q \wedge$$
$$\quad \quad Loves(p, q)$$
$$\quad)$$
$$)$$

Using the predicates

- *Person*(p), which states that p is a person, and
- *Loves*(x, y), which states that x loves y ,

write a sentence in first-order logic that means “there is a person that everyone else loves.”

Answer at

<https://cs103.stanford.edu/pollev>

There is a person that everyone else loves

There is a person p where everyone else loves p

There is a person p where everyone else loves p

“Some A s are B s”

$\exists x. (A(x) \wedge B(x))$

$\exists p. (Person(p) \wedge$
everyone else loves p

)

“Some As are Bs”

$\exists x. (A(x) \wedge B(x))$

$\exists p. (Person(p) \wedge$
everyone else loves p

)

$\exists p. (Person(p) \wedge$
every other person q loves p

)

$\exists p. (Person(p) \wedge$
every person q , other than p , loves p
 $)$

$\exists p. (Person(p) \wedge$
every person q , other than p , loves p

)

“All As are Bs”

$\forall x. (A(x) \rightarrow B(x))$

$\exists p. (Person(p) \wedge$
 $\forall q. (Person(q) \wedge p \neq q \rightarrow$
 $q \text{ loves } p$
)
)

“All As are Bs”

$\forall x. (A(x) \rightarrow B(x))$

$$\begin{aligned} &\exists p. (Person(p) \wedge \\ &\quad \forall q. (Person(q) \wedge p \neq q \rightarrow \\ &\quad \quad q \text{ loves } p) \\ &)\end{aligned}$$

$$\begin{aligned} &\exists p. (Person(p) \wedge \\ &\quad \forall q. (Person(q) \wedge p \neq q \rightarrow \\ &\quad \quad Loves(q, p) \\ &\quad) \\ &) \end{aligned}$$

Quantifier Ordering

Combining Quantifiers

- Most interesting statements in first-order logic require a combination of quantifiers.

“Every person loves someone else”

For every person... $\forall p. (Person(p) \rightarrow$
... there is another person ... $\exists q. (Person(q) \wedge p \neq q \wedge$
... they love $Loves(p, q)$
)
)

Combining Quantifiers

- Most interesting statements in first-order logic require a combination of quantifiers.

“There is someone everyone else loves.”

There is a person...	$\exists p. (Person(p) \wedge$
... that everyone else ...	$\forall q. (Person(q) \wedge p \neq q \rightarrow$
... loves.	$Loves(q, p)$
	)
	)

For Comparison

For every person... $\forall p. (Person(p) \rightarrow$

... there is another person ... $\exists q. (Person(q) \wedge p \neq q \wedge$
... they love $Loves(p, q)$

)
)

There is a person... $\exists p. (Person(p) \wedge$

... that everyone else ... $\forall q. (Person(q) \wedge p \neq q \rightarrow$
... loves. $Loves(q, p)$

)
)

Quantifier Ordering

- Consider these two first-order formulas:

$$\forall m. \exists n. m < n.$$

$$\exists n. \forall m. m < n.$$

- Pretend for the moment that our world consists purely of natural numbers, so the variables m and n refer specifically to natural numbers.
- One of these statements is true. The other is false.
- Which is which?
- Why?

Answer at

<https://cs103.stanford.edu/pollev>

Quantifier Ordering

- Consider these two first-order formulas:

$$\forall m. \exists n. m < n.$$

$$\exists n. \forall m. m < n.$$

- This says

**for every natural number m ,
there's a larger natural number n .**

- This is true: given any $m \in \mathbb{N}$, we can choose n to be $m + 1$.
- Notice that we can pick n based on m , and we don't have to pick the same n each time.

Quantifier Ordering

- Consider these two first-order formulas:

$$\forall m. \exists n. m < n.$$

$$\exists n. \forall m. m < n.$$

- This says

**there is a natural number n
that's larger than every natural number m**

- This is false: no natural number is bigger than every natural number.
- Because $\exists n$ comes first, we have to make a single choice of n that works regardless of what we choose for m .

Quantifier Ordering

- The statement

$$\forall x. \exists y. P(x, y)$$

means “for any choice of x , there's some choice of y where $P(x, y)$ is true.”

- The choice of y can be different every time and can depend on x .

Quantifier Ordering

- The statement

$$\exists x. \forall y. P(x, y)$$

means “there is some x where for any choice of y , we get that $P(x, y)$ is true.”

- Since the inner part has to work for any choice of y , this places a lot of constraints on what x can be.

Order matters when mixing existential
and universal quantifiers!

Time-Out for Announcements!

Problem Set Two

- Problem Set One was due today at 1:00PM.
 - You can extend the deadline to 1:00PM Saturday using one of your late days. As usual, no late submissions will be accepted beyond 1:00PM Saturday without prior approval.
- Problem Set Two goes out today. It's due next Friday at 1:00PM.
 - Explore first-order logic!
 - Expand your proofwriting toolkit!
- We have some online readings for this problem set.
 - Check out the ***Guide to Logic Translations*** for more on how to convert from English to FOL.
 - Check out the ***Guide to Negations*** for information about how to negate formulas.
 - Check out the ***First-Order Translation Checklist*** for details on how to check your work.

Reminder: Stanford Honor Code

- As a reminder on course policies:
 - ChatGPT and other generative AI tools are off-limits for graded work.
 - You can discuss high-level ideas with other students, but can only share concrete solutions with your problem set partner.
- If you submitted something you shouldn't have, keep an eye out for the Regret Clause Form that will go out this weekend.
- ***We take the Honor Code seriously.*** It promotes learning and basic fairness.

A Music Recommendation



Back to CS103!

Mechanics: Negating Statements

An Extremely Important Table

	When is this true?	When is this false?
$\forall x. P(x)$	For all objects x , $P(x)$ is true.	There is an x where $P(x)$ is false.
$\exists x. P(x)$	There is an x where $P(x)$ is true.	For all objects x , $P(x)$ is false.
$\forall x. \neg P(x)$	For all objects x , $P(x)$ is false.	There is an x where $P(x)$ is true.
$\exists x. \neg P(x)$	There is an x where $P(x)$ is false.	For all objects x , $P(x)$ is true.

An Extremely Important Table

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An Extremely Important Table

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$\exists x. \neg P(x)$	There is an x where $P(x)$ is false.	$\forall x. P(x)$

An Extremely Important Table

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$\exists x. \neg P(x)$	There is an x where $P(x)$ is false.	$\forall x. P(x)$

Negating First-Order Statements

- Use the equivalences

$\neg \forall x. A$ is equivalent to $\exists x. \neg A$

$\neg \exists x. A$ is equivalent to $\forall x. \neg A$

to negate quantifiers.

- Mechanically:
 - Push the negation across the quantifier.
 - Change the quantifier from $\forall$ to $\exists$ or vice-versa.
- Use techniques from propositional logic to negate connectives.

Taking a Negation

$\forall x. \exists y. \text{Loves}(x, y)$
(*“Everyone loves someone.”*)

$\neg \forall x. \exists y. \text{Loves}(x, y)$
 $\exists x. \neg \exists y. \text{Loves}(x, y)$
 $\exists x. \forall y. \neg \text{Loves}(x, y)$
(*“There's someone who doesn't love anyone.”*)

Two Useful Equivalences

- The following equivalences are useful when negating statements in first-order logic:

$\neg(p \wedge q)$ *is equivalent to* $p \rightarrow \neg q$

$\neg(p \rightarrow q)$ *is equivalent to* $p \wedge \neg q$

- These identities are useful when negating statements involving quantifiers.
 - $\wedge$ is used in existentially-quantified statements.
 - $\rightarrow$ is used in universally-quantified statements.
- When pushing negations across quantifiers, we ***strongly recommend*** using the above equivalences to keep $\rightarrow$ with $\forall$ and $\wedge$ with $\exists$.

Negating Quantifiers

- What is the negation of the following statement, which says “there is a cute puppy”?

$\exists x. (Puppy(x) \wedge Cute(x))$

Answer at

<https://cs103.stanford.edu/pollev>

Negating Quantifiers

- What is the negation of the following statement, which says “there is a cute puppy”?

$$\exists x. (\textit{Puppy}(x) \wedge \textit{Cute}(x))$$

- We can obtain it as follows:

$$\neg \exists x. (\textit{Puppy}(x) \wedge \textit{Cute}(x))$$

$$\forall x. \neg (\textit{Puppy}(x) \wedge \textit{Cute}(x))$$

$$\forall x. (\textit{Puppy}(x) \rightarrow \neg \textit{Cute}(x))$$

- This says “no puppy is cute.”
- Do you see why this is the negation of the original statement from both an intuitive and formal perspective?

$$\exists S. (Set(S) \wedge \forall x. x \notin S)$$

("There is a set with no elements.")

$$\neg \exists S. (Set(S) \wedge \forall x. x \notin S)$$

$$\forall S. \neg (Set(S) \wedge \forall x. \neg x \notin S)$$

$$\forall S. (Set(S) \rightarrow \neg \forall x. x \notin S)$$

$$\forall S. (Set(S) \rightarrow \exists x. \neg (x \notin S))$$

$$\forall S. (Set(S) \rightarrow \exists x. x \in S)$$

("Every set contains at least one element.")

Restricted Quantifiers

Quantifying Over Sets

- The notation

$$\forall x \in S. P(x)$$

means “for any element x of set S , $P(x)$ holds.” (It’s vacuously true if S is empty.)

- The notation

$$\exists x \in S. P(x)$$

means “there is an element x of set S where $P(x)$ holds.” (It’s false if S is empty.)

Quantifying Over Sets

- The syntax

$$\forall x \in S. P(x)$$

$$\exists x \in S. P(x)$$

is allowed for quantifying over sets.

- In CS103, feel free to use these restricted quantifiers, but please do not use variants of this syntax.
- For example, don't do things like this:

$$\forall x \text{ with } P(x). Q(x)$$

$$\forall y \text{ such that } P(y) \wedge Q(y). R(y).$$

$$\exists P(x). Q(x)$$

Expressing Uniqueness

Using the predicate

- *WayToFindOut*(w), which states that w is a way to find out,

write a sentence in first-order logic that means “there is only one way to find out.”

There is only one way to find out.

Something is a way to find out, and nothing else is.

Some thing w is a way to find out, and nothing else is.

*Some thing w is a way to find out, and nothing besides w
is a way to find out*

$\exists w. (WayToFindOut(w) \wedge$
 nothing besides w is way to find out
)

$\exists w. (WayToFindOut(w) \wedge$
 anything that isn't w isn't a way to find out
 $)$

$\exists w. (WayToFindOut(w) \wedge$
 any thing x that isn't w isn't a way to find out
 $)$

$\exists w. (WayToFindOut(w) \wedge$
 $\forall x. (x \neq w \rightarrow x \text{ isn't a way to find out})$
 $)$

$$\exists w. (WayToFindOut(w) \wedge$$
$$\quad \forall x. (x \neq w \rightarrow \neg WayToFindOut(x))$$
$$)$$

$$\exists w. (WayToFindOut(w) \wedge$$
$$\quad \forall x. (x \neq w \rightarrow \neg WayToFindOut(x))$$
$$)$$

$$\exists w. (WayToFindOut(w) \wedge \\ \forall x. (WayToFindOut(x) \rightarrow x = w) \\)$$

$$\exists w. (WayToFindOut(w) \wedge \\ \quad \forall x. (WayToFindOut(x) \rightarrow x = w) \\)$$

Expressing Uniqueness

- To express the idea that there is exactly one object with some property, we write that
 - there exists at least one object with that property, and that
 - there are no other objects with that property.
- You sometimes see a special “uniqueness quantifier” used to express this:

$$\exists!x. P(x)$$

- For the purposes of CS103, please do not use this quantifier. We want to give you more practice using the regular $\forall$ and $\exists$ quantifiers.

Next Time

- ***Functions***
 - How do we model transformations and pairings?
- ***First-Order Definitions***
 - Where does first-order logic come into all of this?
- ***Proofs with Definitions***
 - How does first-order logic interact with proofs?